



Position and Orientation in Ad Hoc Networks

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summary

- problem and motivation
- example
- basic idea
- AoA capable nodes
- algorithm outline
 - bearing propagation (DV)
 - error control
- simulation results
- conclusions

problem statement

ad hoc deployed nodes should be able to know their

○ location

- global coordinates
- low overhead for mobility
- accuracy comparable with the node communication range
- independent operation for disconnected regions
- without additional infrastructure

○ orientation

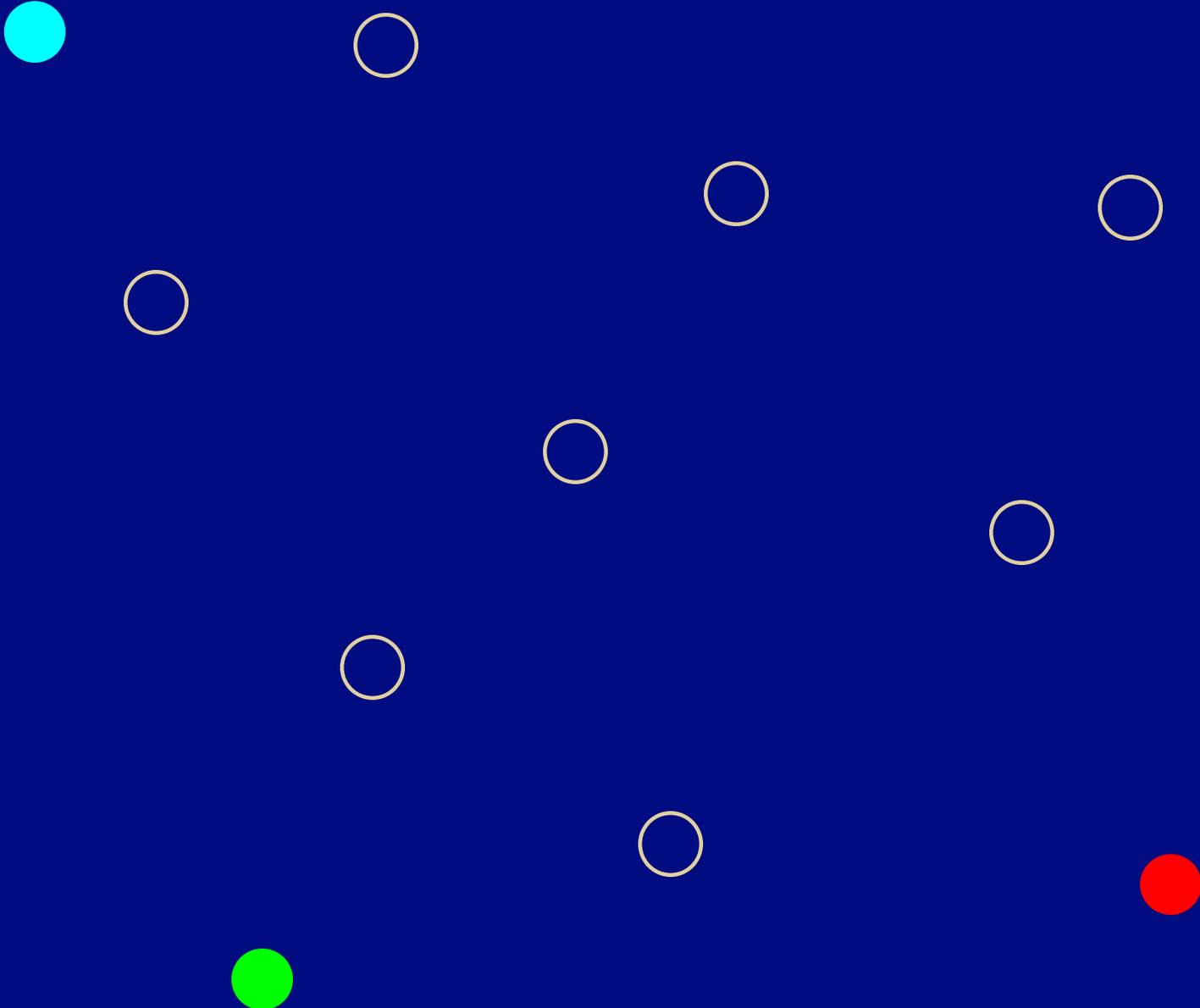
- heading

motivation

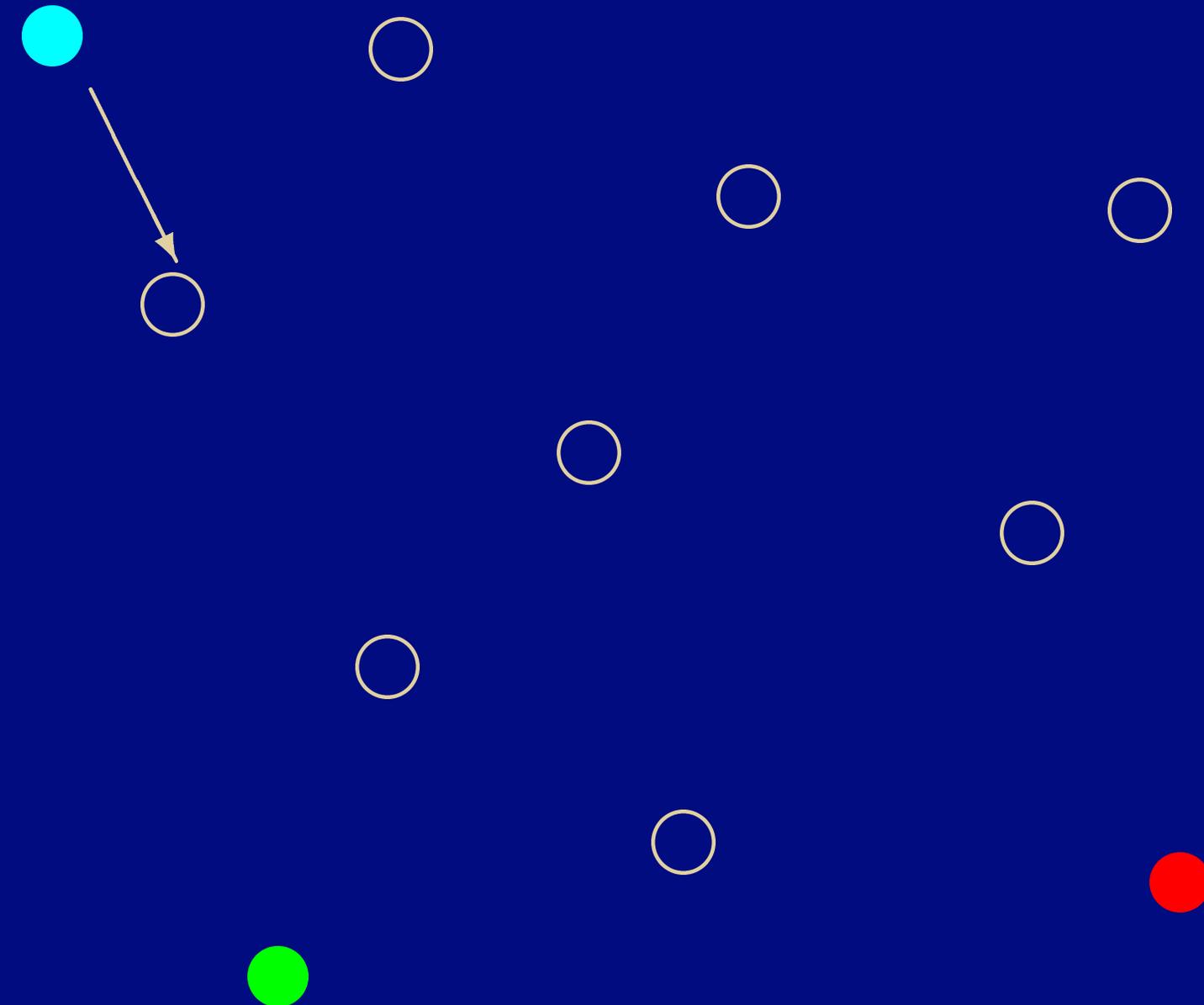
- a sensor reports a phenomenon and its:
 - position
 - place it on a map
 - routing with small or no routing tables
 - orientation
 - remote navigation
 - fine grained control – camera orientation
 - intensity

- possible solutions
 - GPS + digital compass in each node
 - compasses do not work well indoors
 - GPS needs line of sight

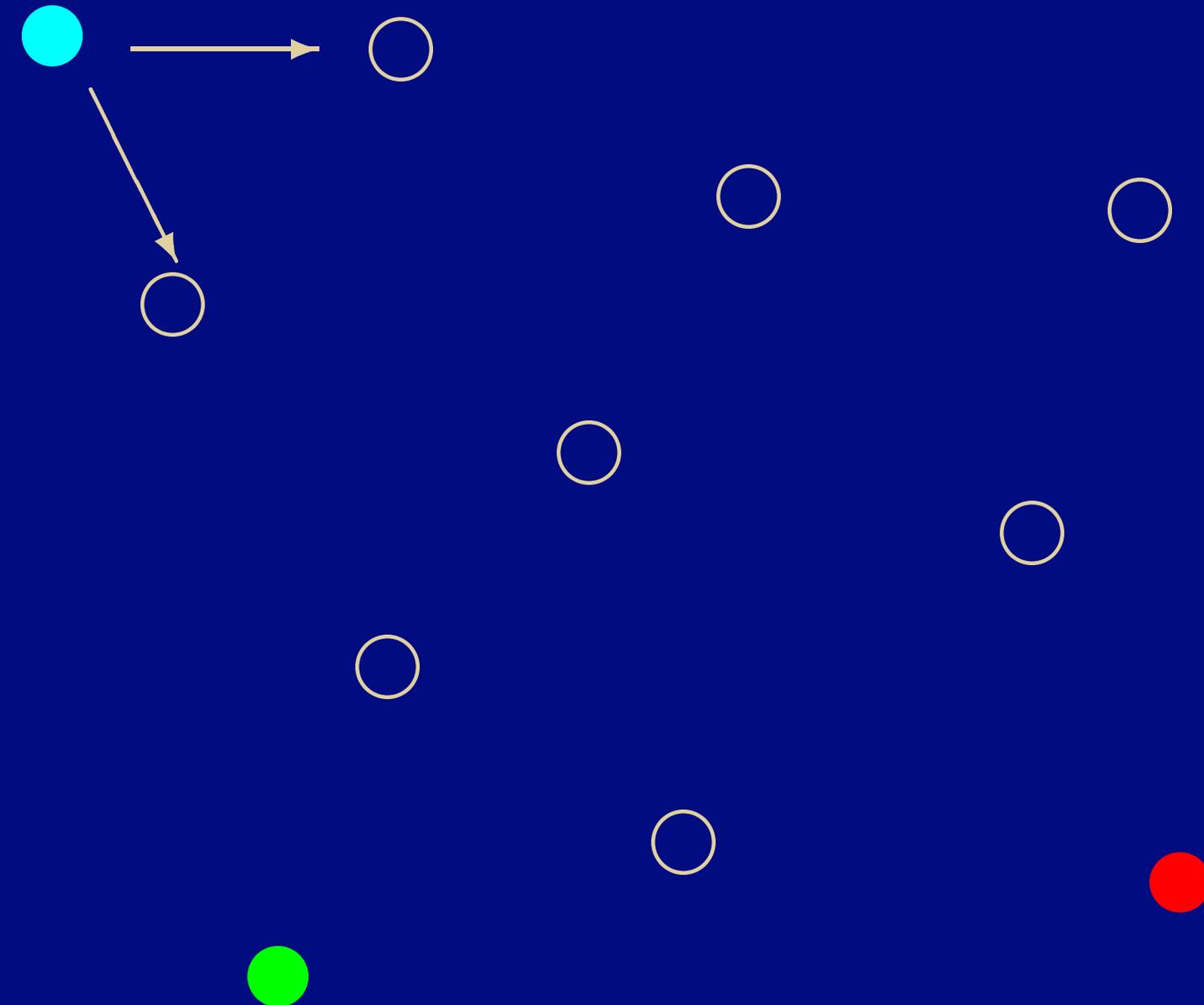
example



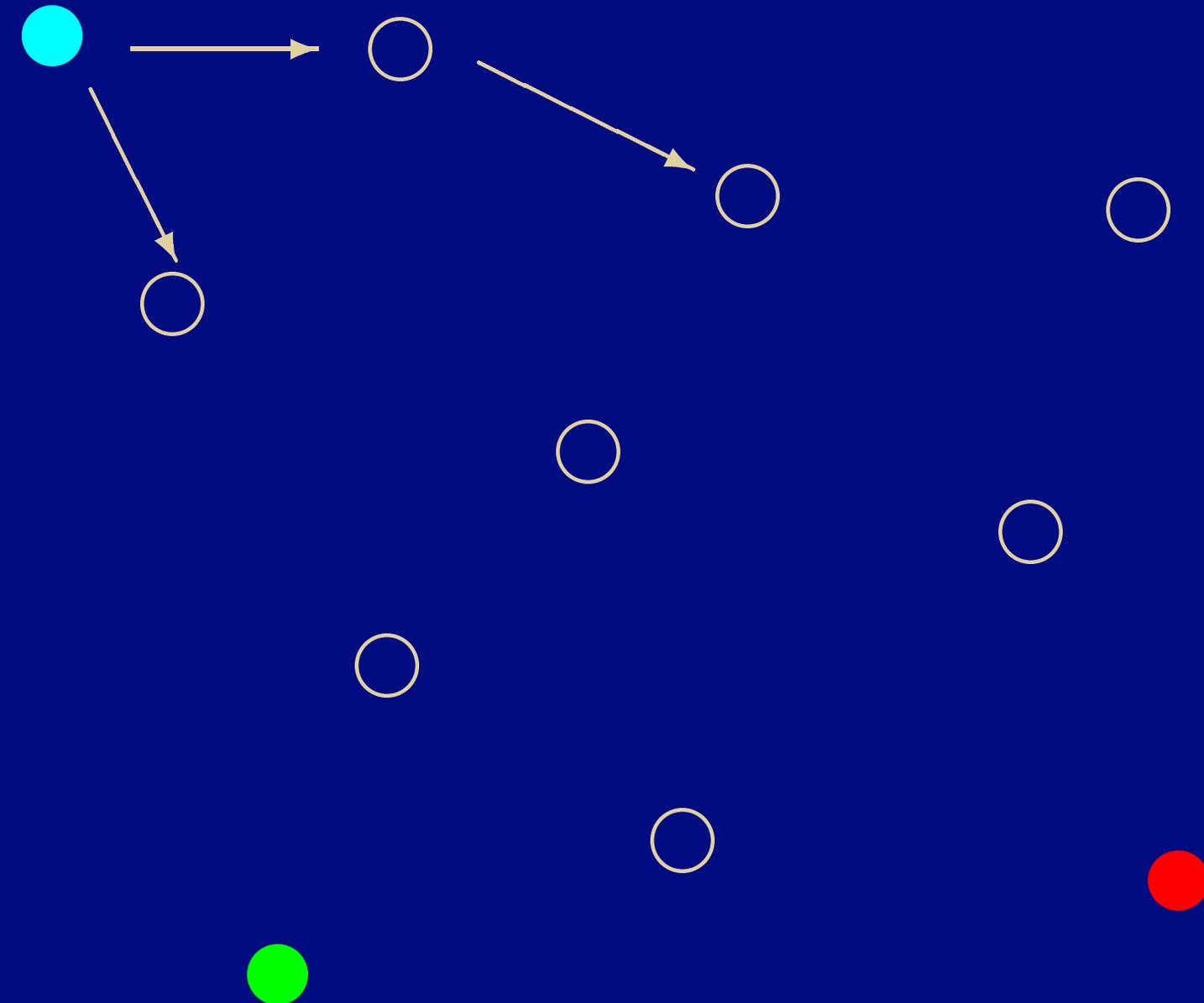
example



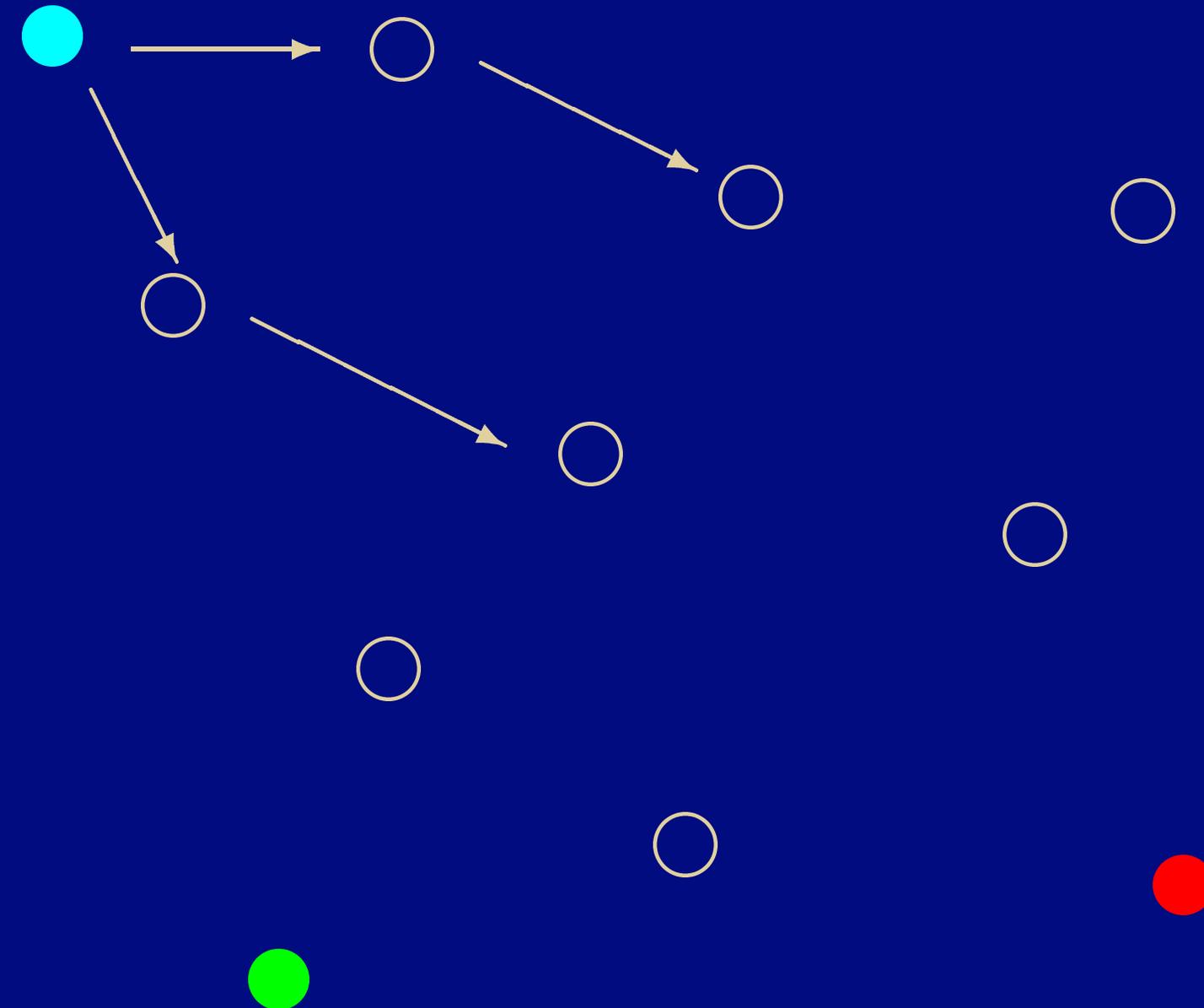
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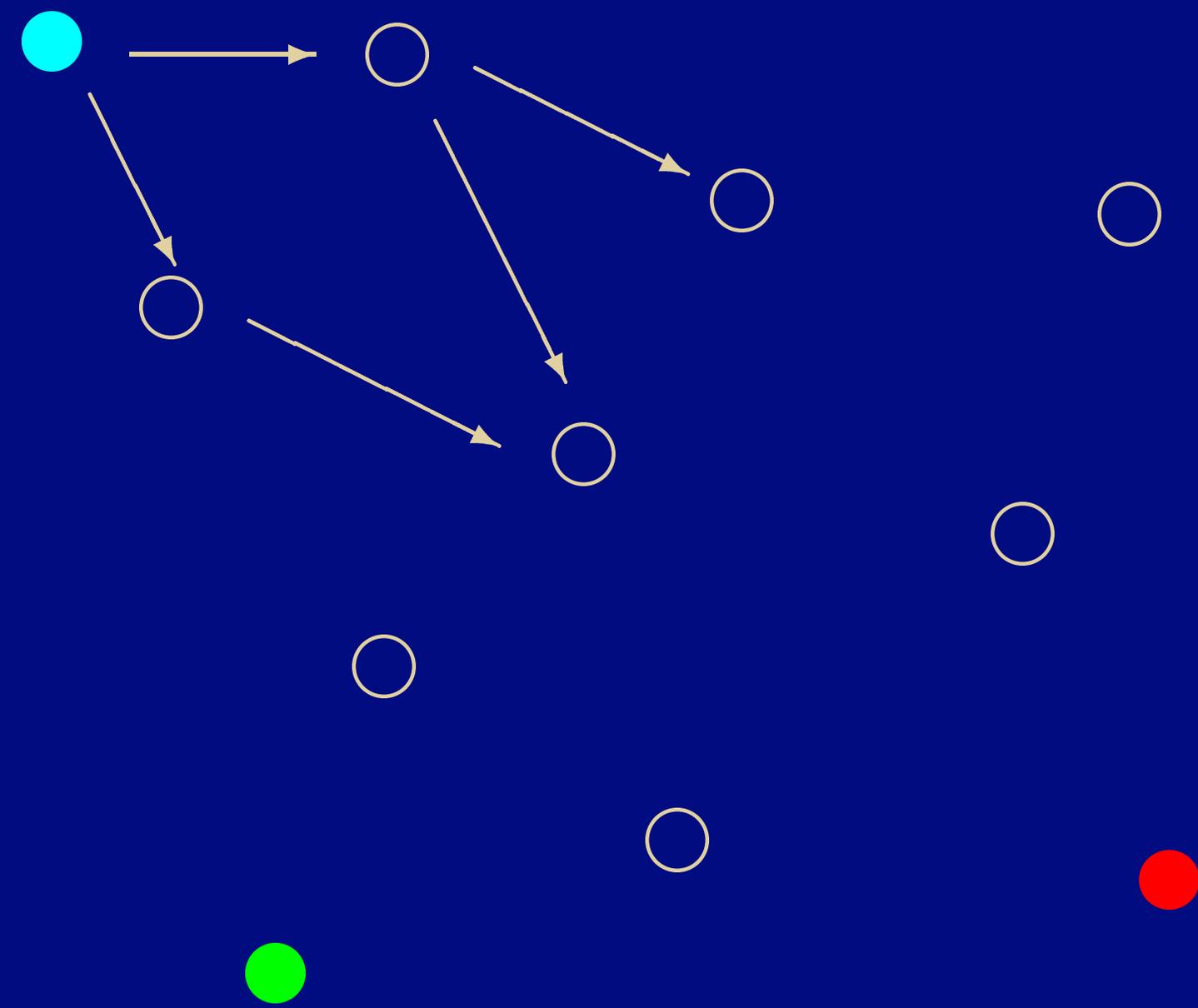
example ✂



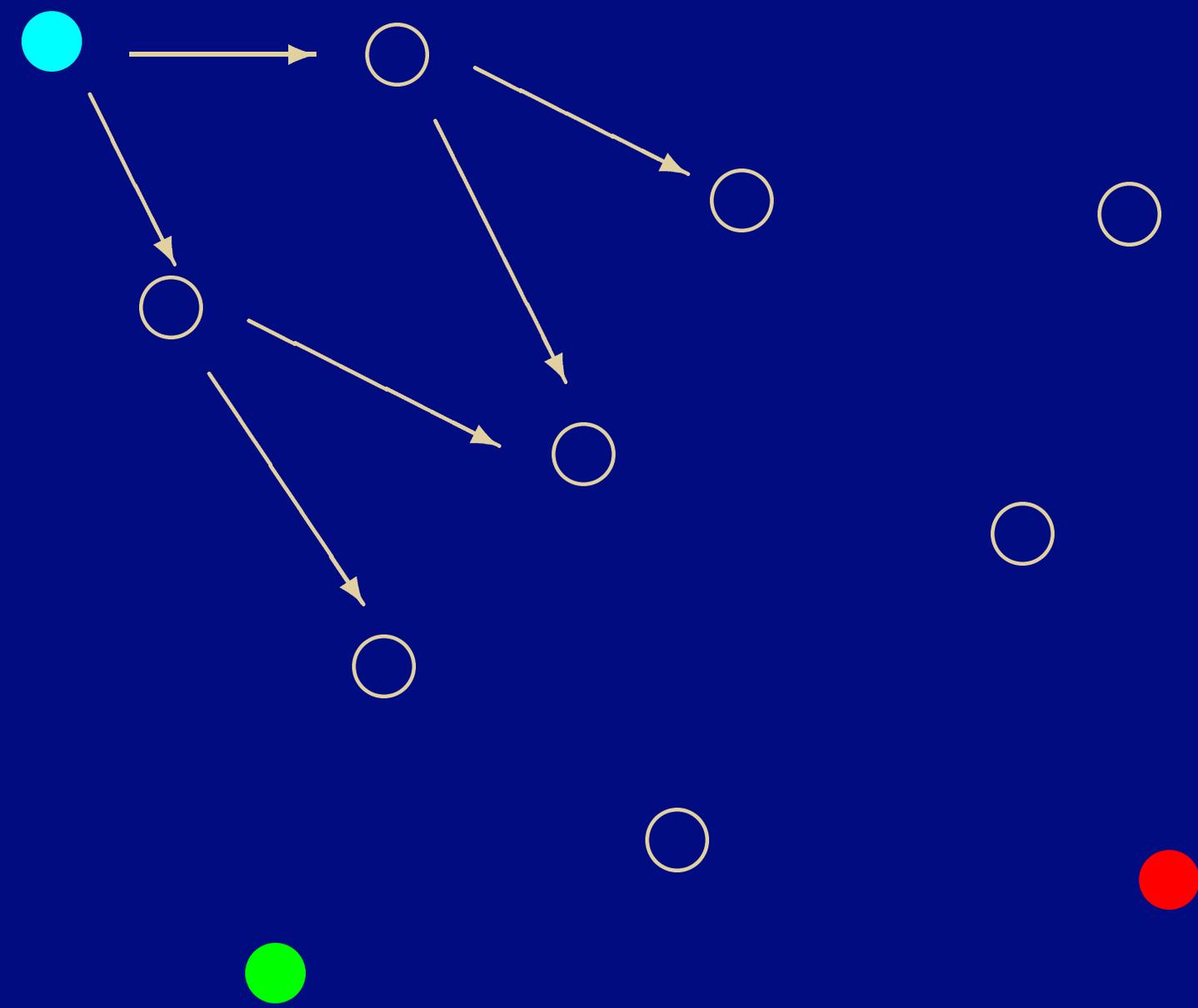
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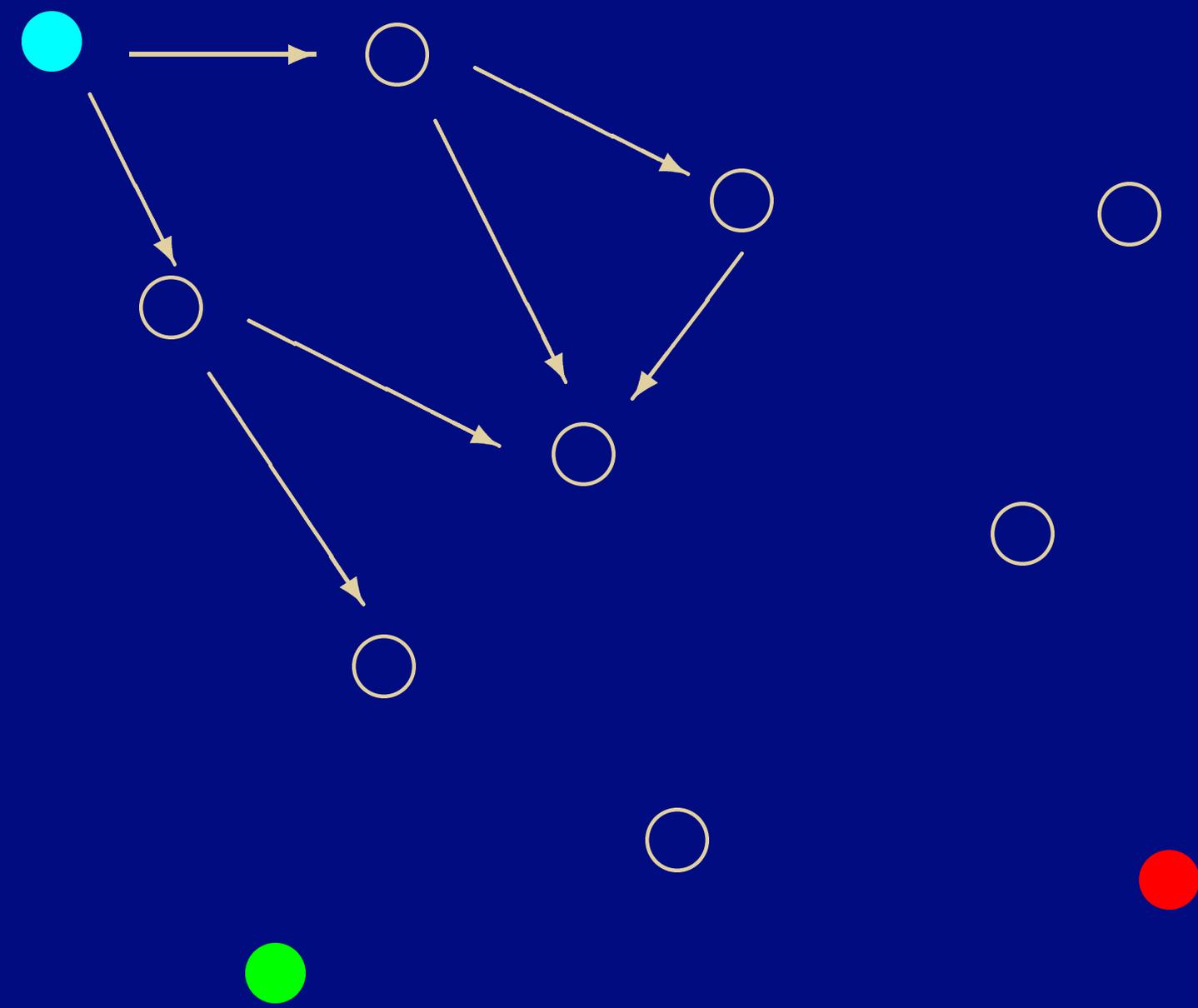
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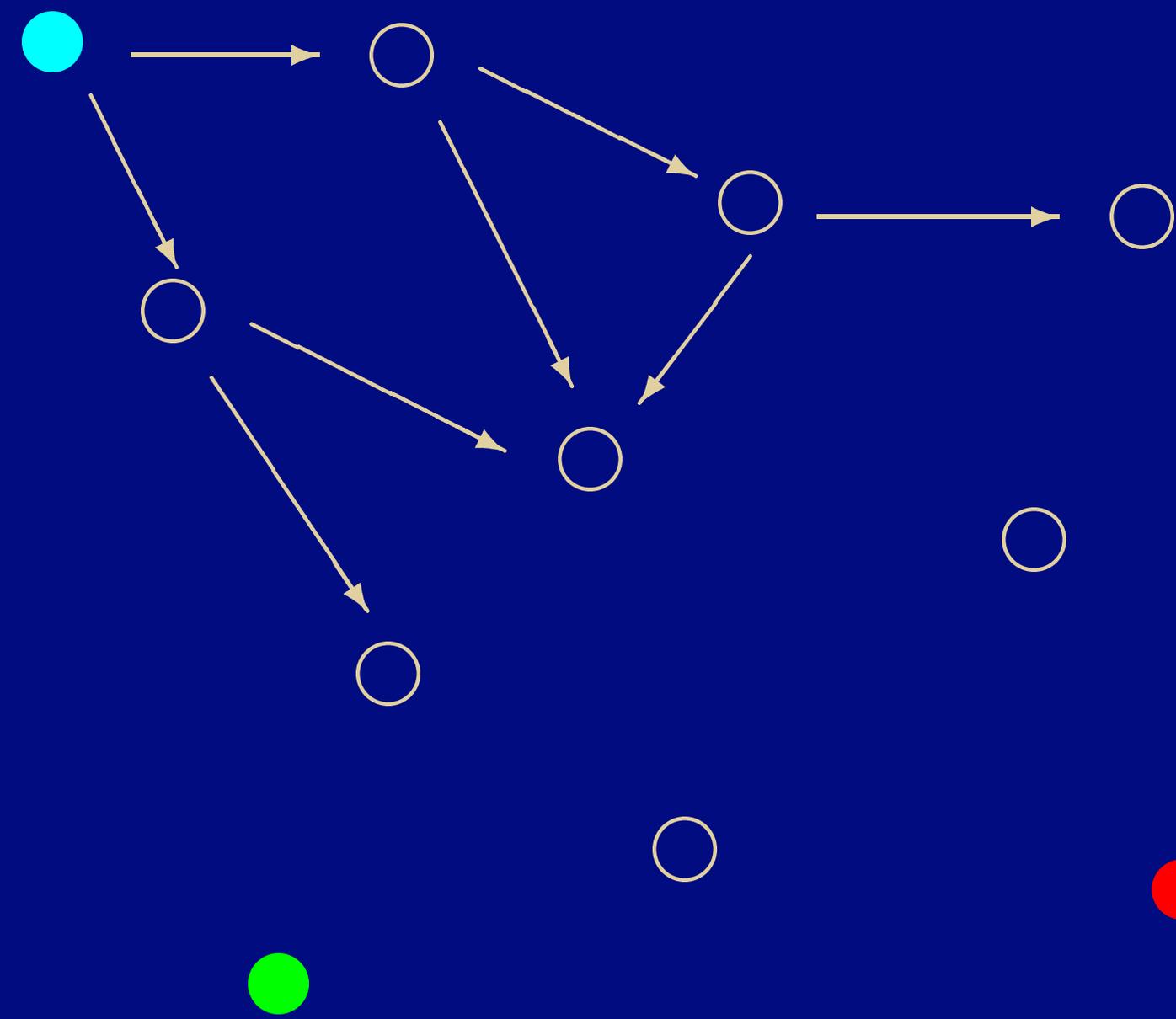
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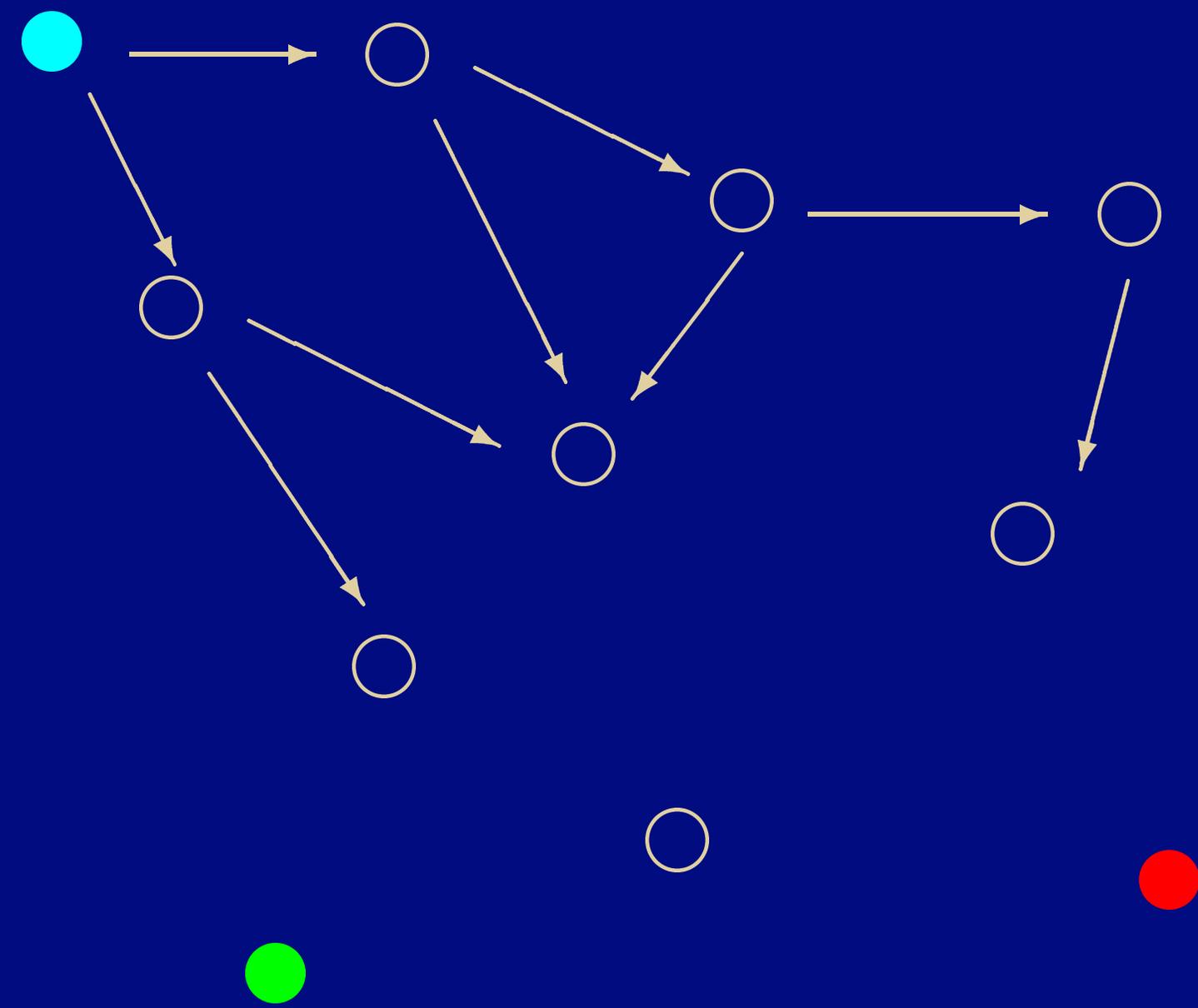
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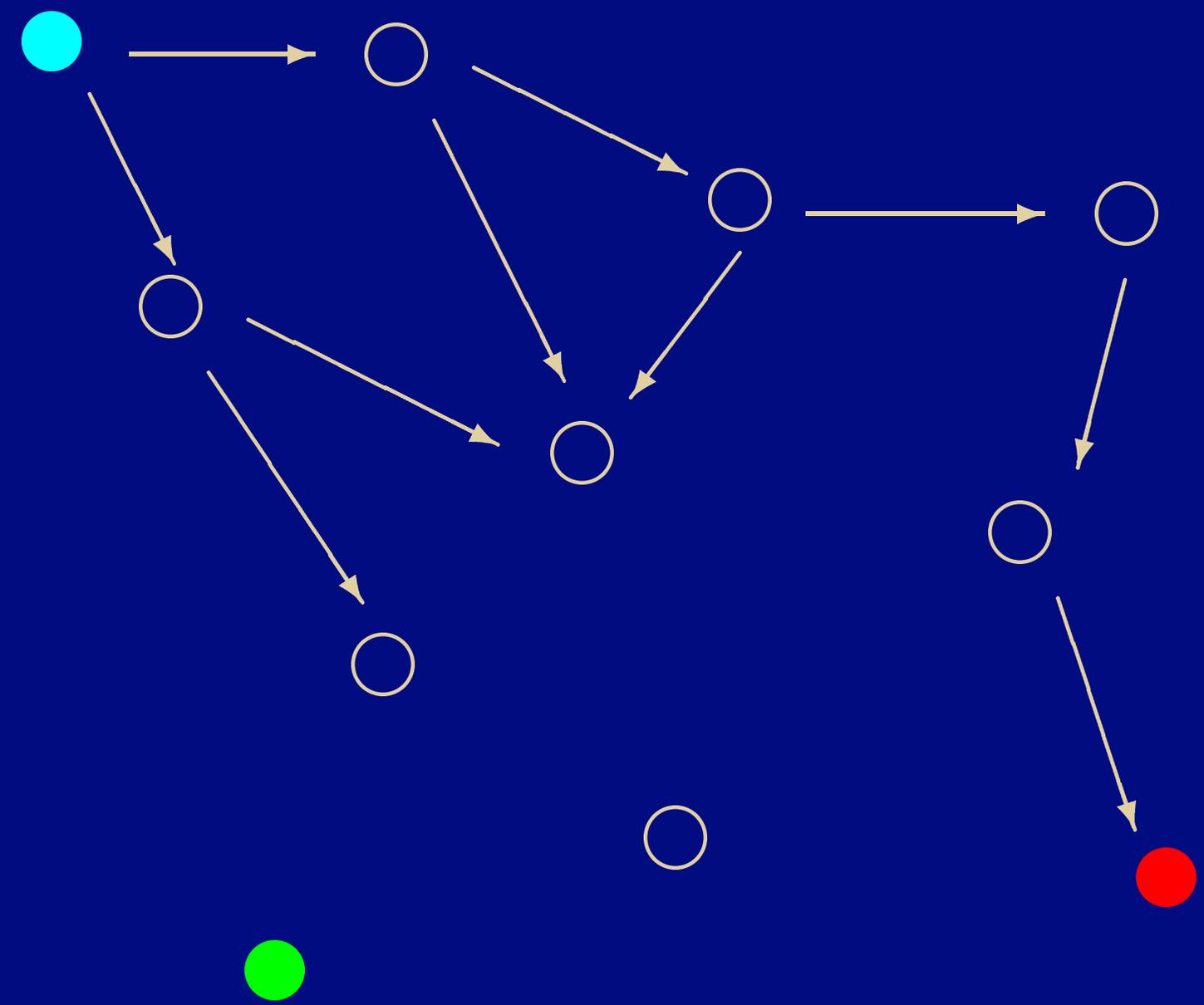
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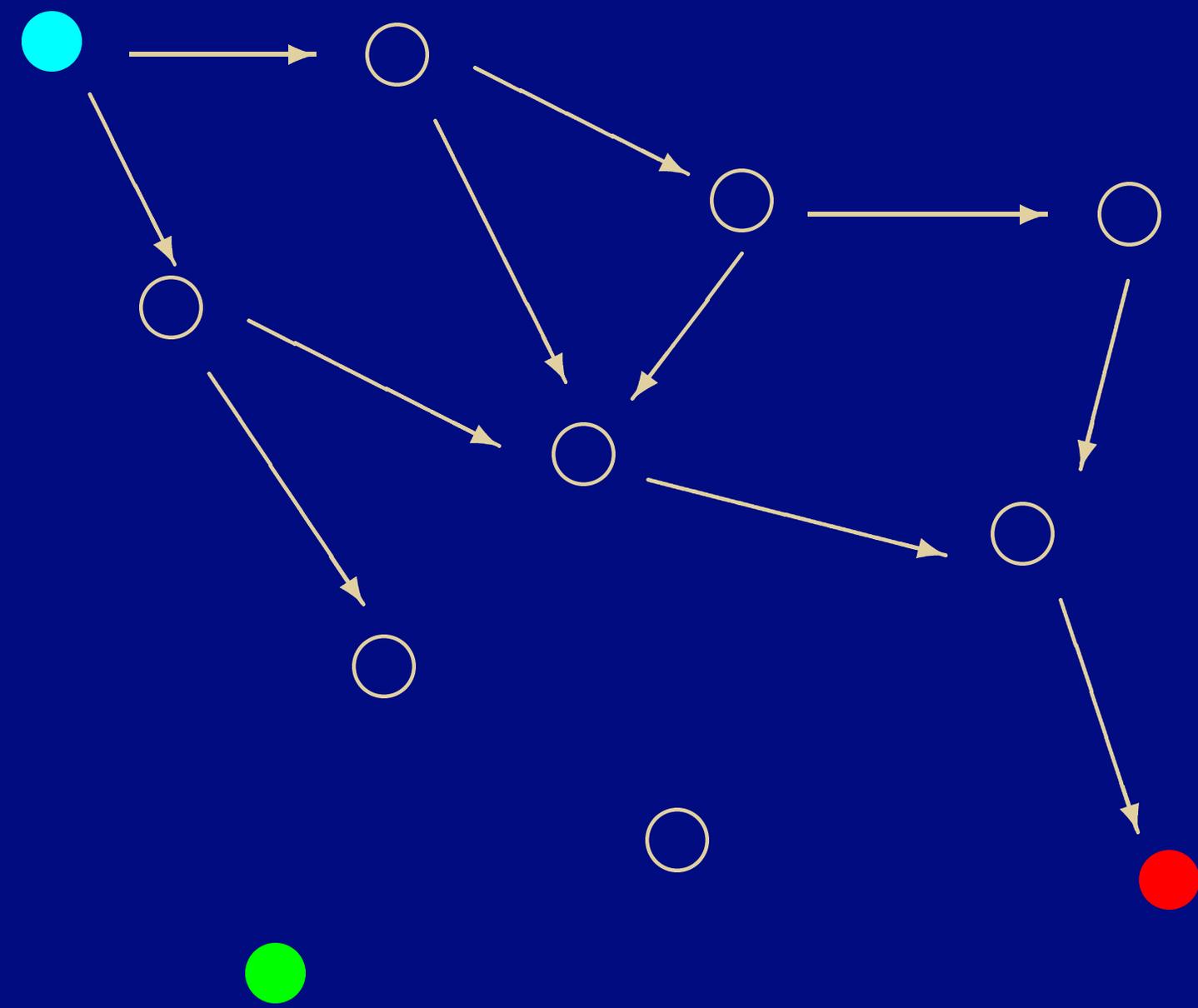
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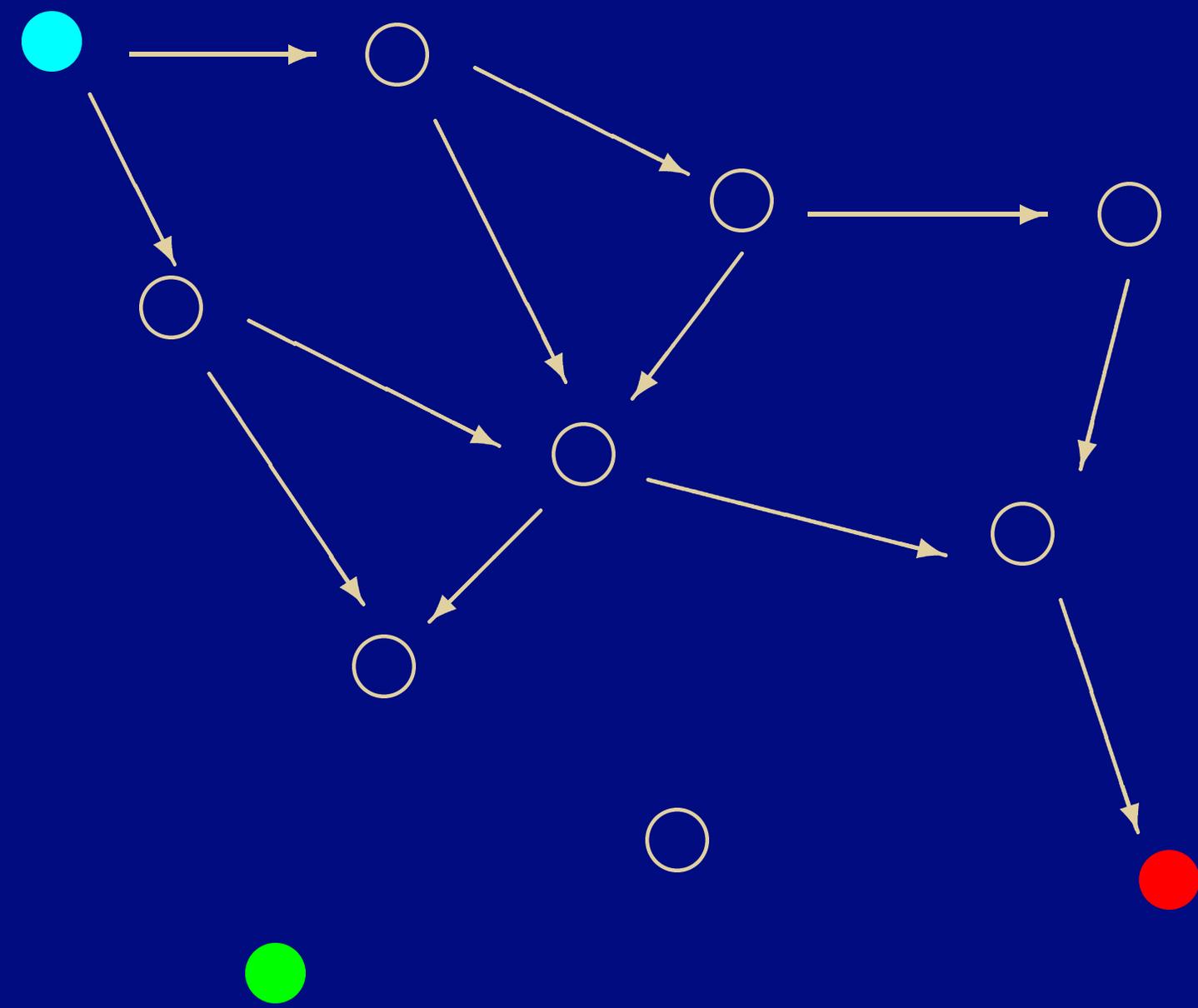
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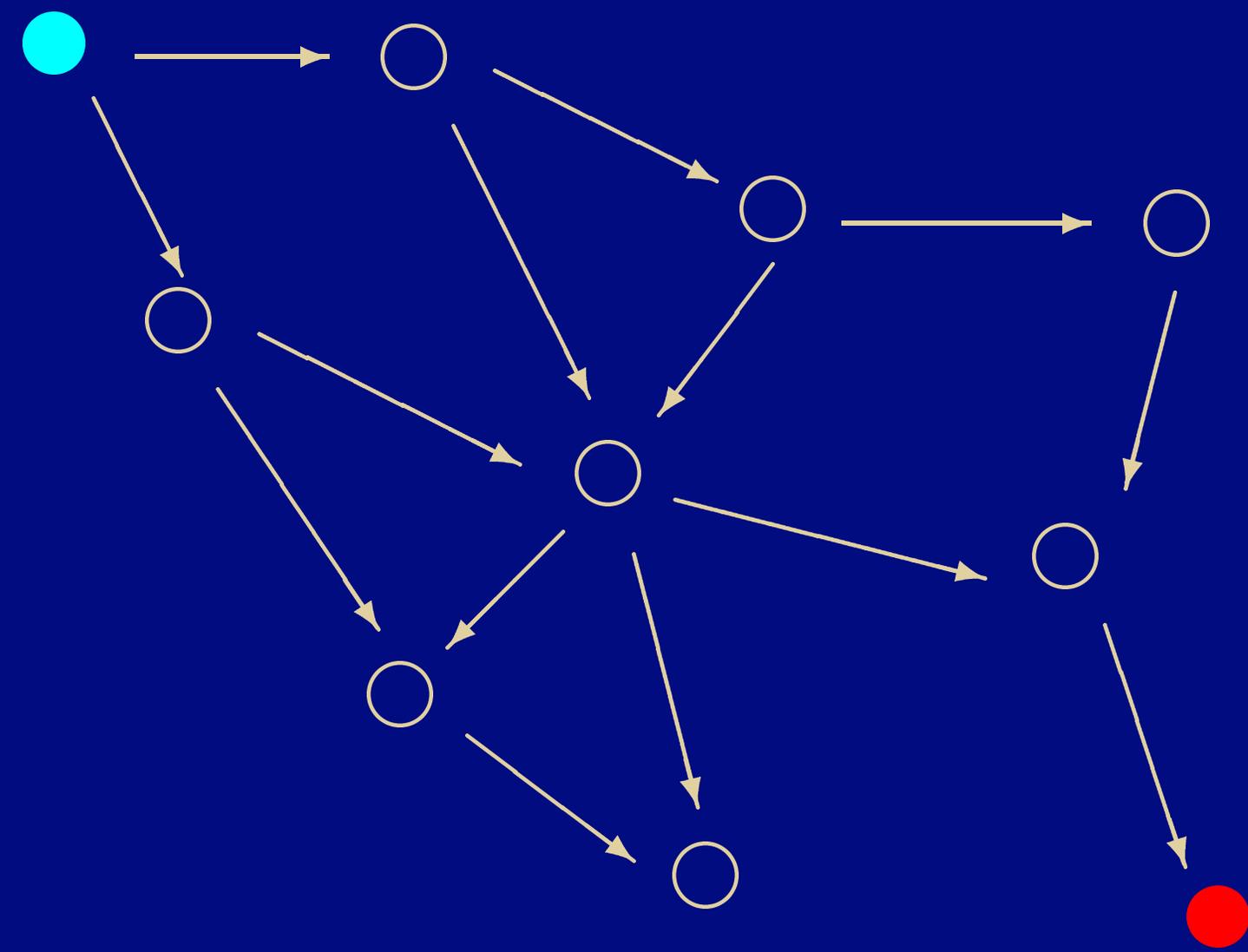
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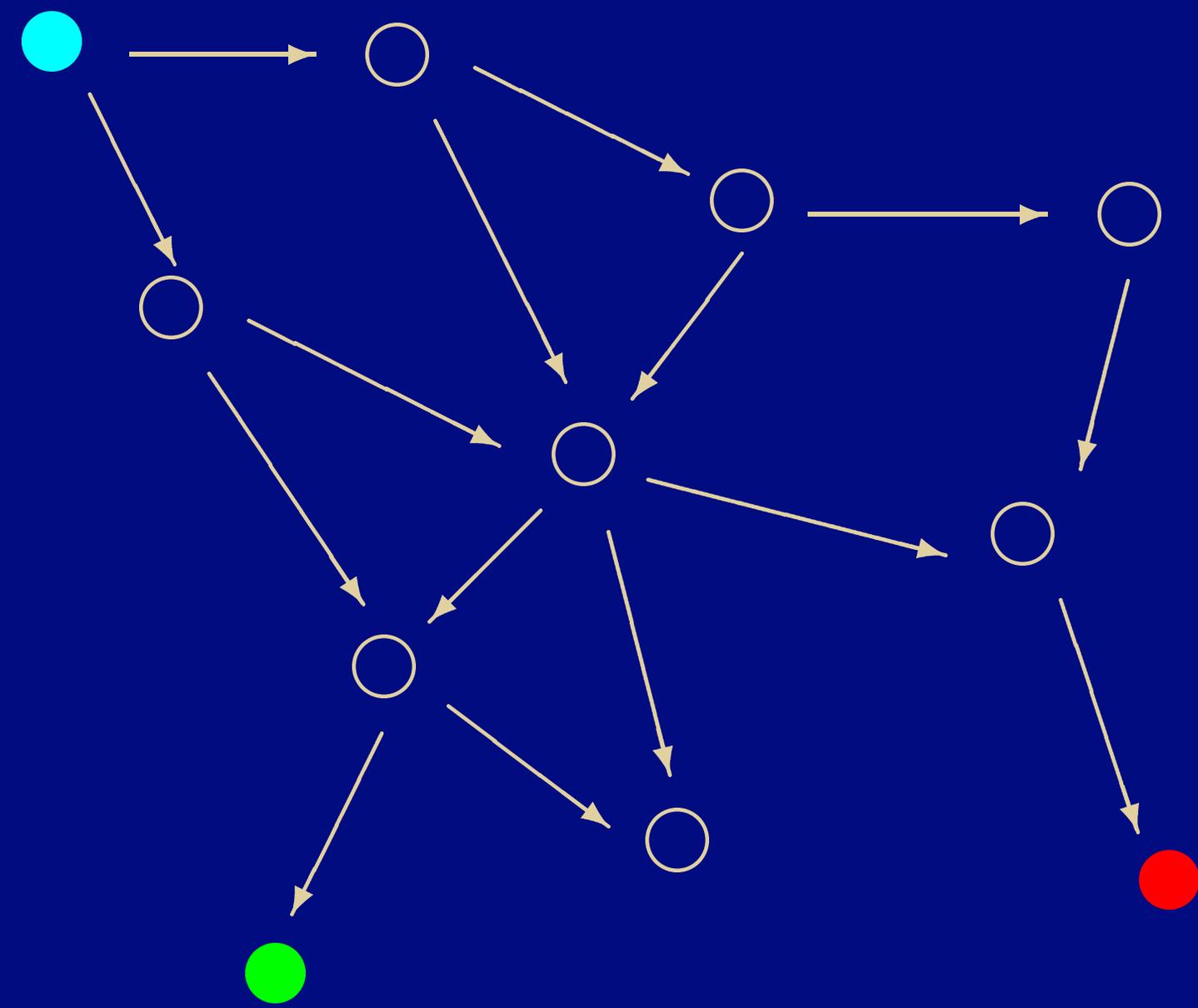
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example ✂



example ✂



terminology and assumptions

○ terminology

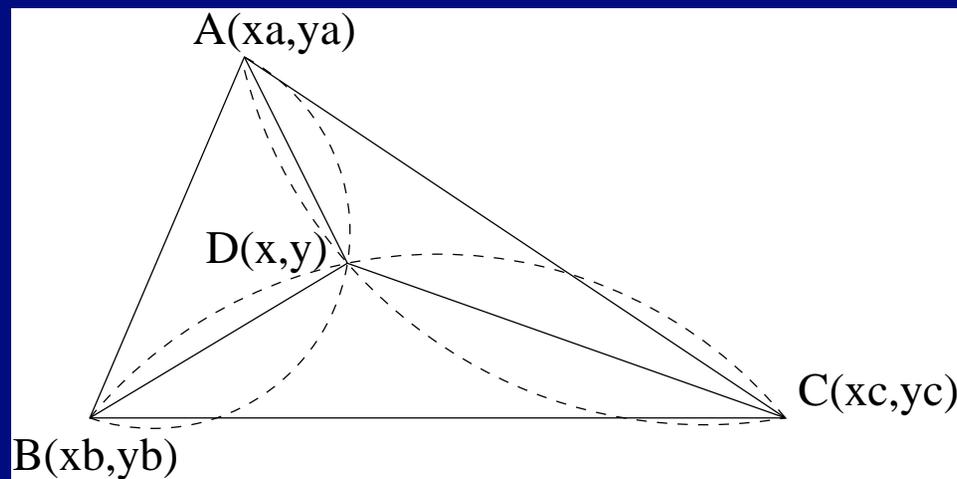
- landmark - a node which knows its own position
- range - distance to an objective
- bearing - angle between one's facing direction and some objective
- heading, orientation - absolute bearing, or angle to north

○ assumptions

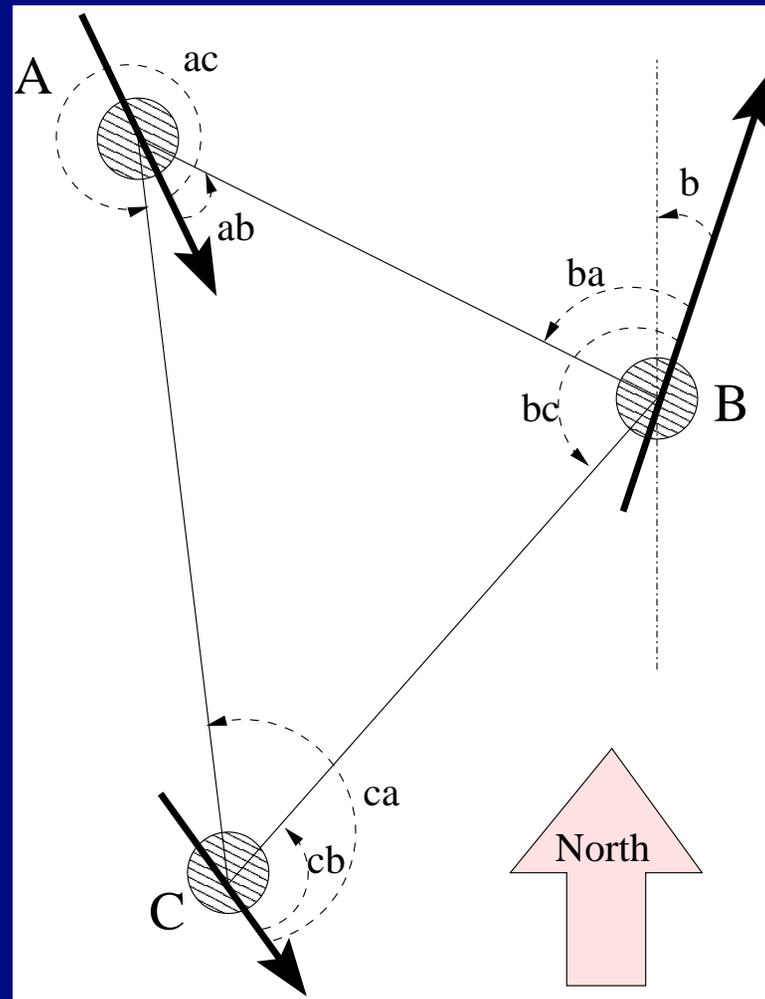
- nodes are deployed randomly → connected graph
- a fixed fraction of nodes are landmarks
- nodes have a limited radio range

basic idea

- Given
 - (imprecise) bearings to at least three landmarks , \widehat{ADB} , \widehat{BDC} , \widehat{CDA}
 - the positions of the landmarks (x_i, y_i) , $i = A, B, C$
- a node may infer its own position (x, y)
- and its orientation/heading, once (x, y) and (x_A, y_A) are known



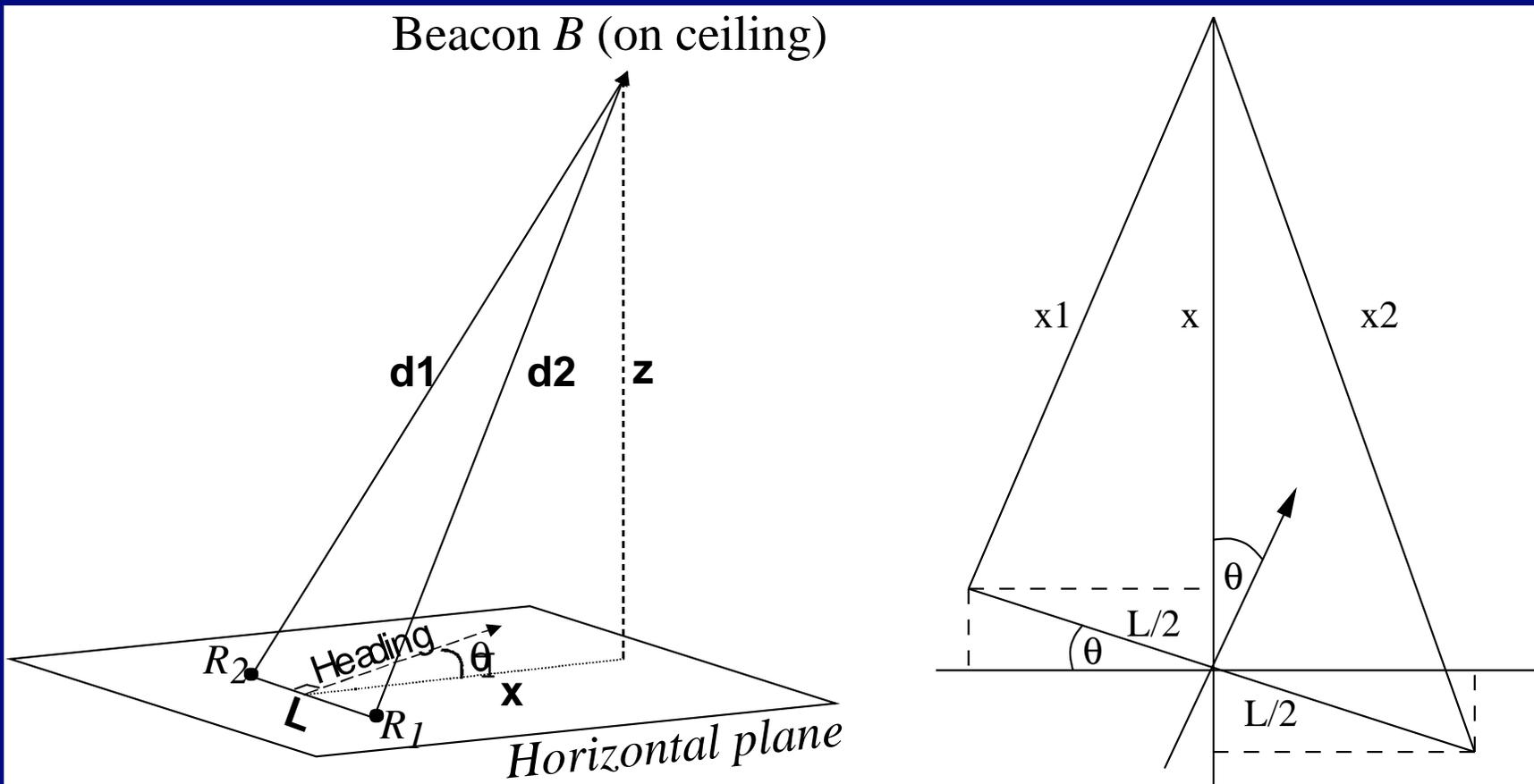
angle of arrival capability - example



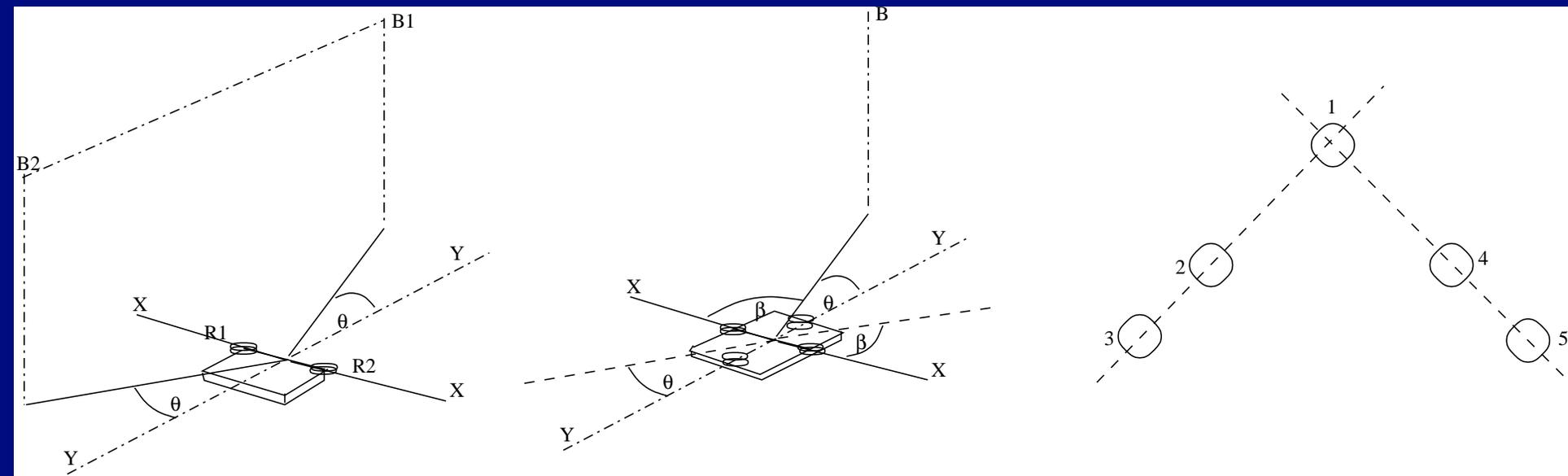
angle of arrival capable nodes

- Cricket compass - MIT (Hari Balakrishnan, Mobicom 2001)
 - uses 5 ultrasound receivers
 - 0.8cm each
 - a few centimeters across
 - TDoA (time difference of arrival)
 - $\pm 10\%$ accuracy for angles < 40 degrees
- Medusa node - UCLA (Mani Srivastava, Mobicom 2001)
- antenna arrays

“Cricket” compass - basic principle



“Cricket” compass - disambiguation



- based on measuring ranges
 - front - back ambiguity
- range differentials are measured using phase difference
 - wavelength $\sim L \rightarrow$ range ambiguity

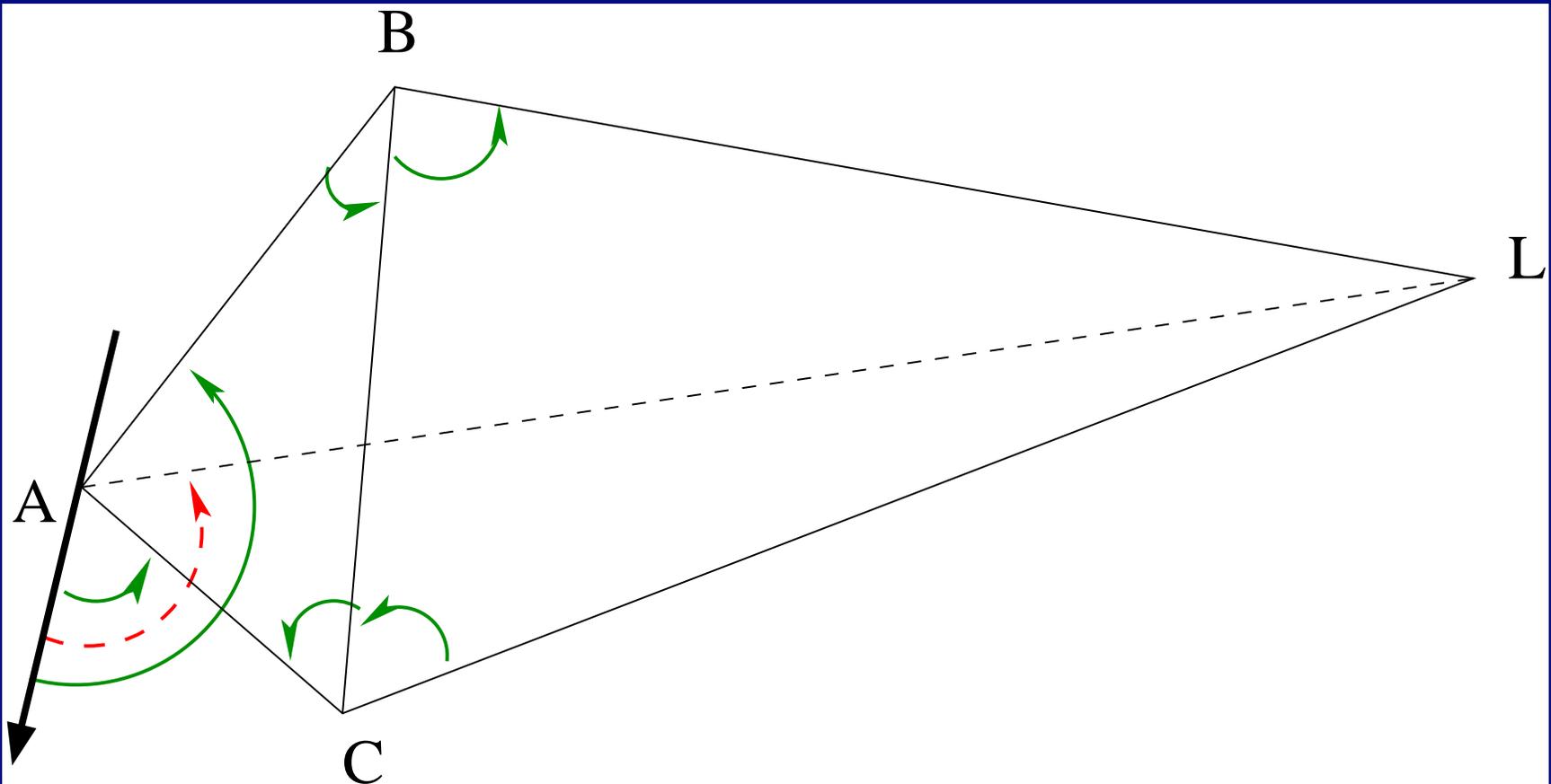
“Medusa” node



algorithm outline

- a few nodes (landmarks) know their position
- all nodes have the AoA capability → find bearings to immediate neighbors
- regular nodes infer bearings to at least three landmarks
 - non-collinear
 - non-cocircular (with the node itself)
- like in DV, bearings to landmarks are propagated hop by hop
- each landmark is treated independently at each node

bearing propagation (DV)



○ green angles are known → find bearing to L (red angle)

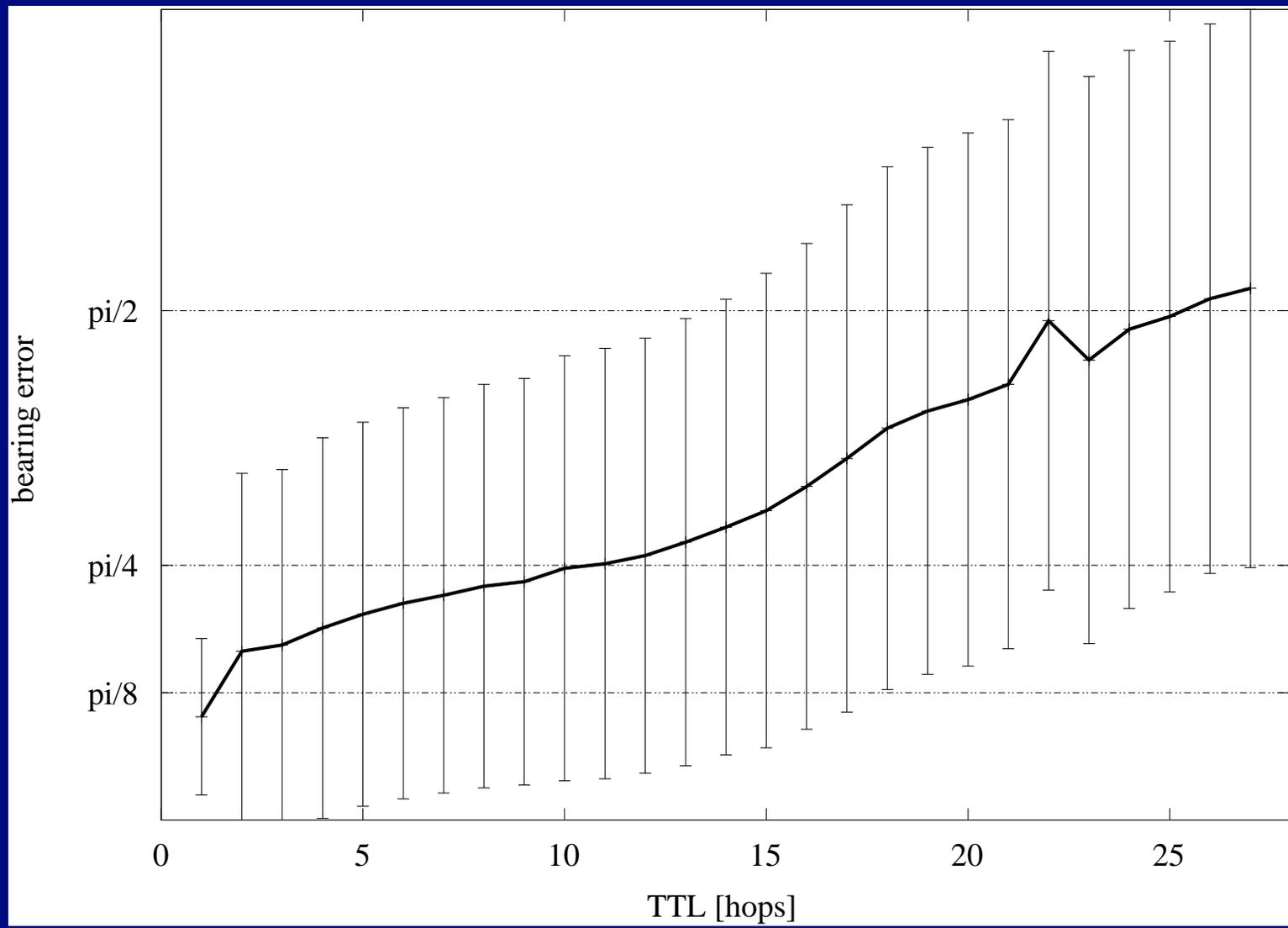
position computation

- each node obtains a table $\{X_i, Y_i, dir_i\}$ - coordinates and bearings to landmarks
- several ways to solve:
 - solve a nonlinear system to intersect the n circles
 - $\binom{n}{2}$ pairs of landmarks, find distances to centres \rightarrow GPS problem
 - $\binom{n}{3}$ triplets of landmarks, find $\binom{n}{3}$ estimates \rightarrow centroid
 - there are $O(n)$ algorithms providing the same accuracy
- absolute position + bearing to known point = heading(absolute orientation)

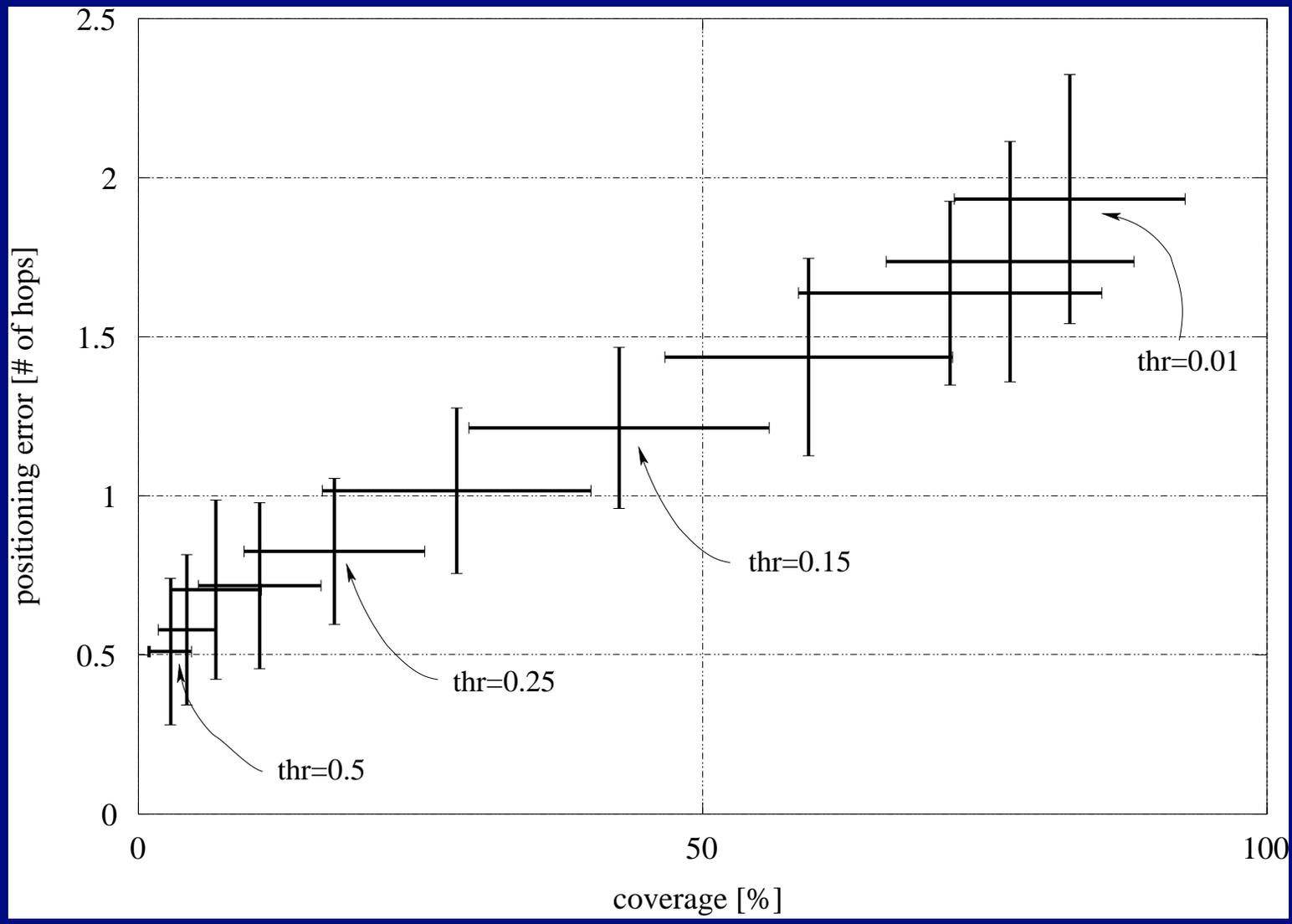
error control

- lightweight methods → CPU, memory, communication
- the propagation scheme compounds errors
 - limit the travel distance of packets (TTL)
- small angle error → large distance error
 - avoid angles below a threshold
- large errors are clustered
 - prune position estimations that are far off the centroid
- these three methods together reduce the errors by half

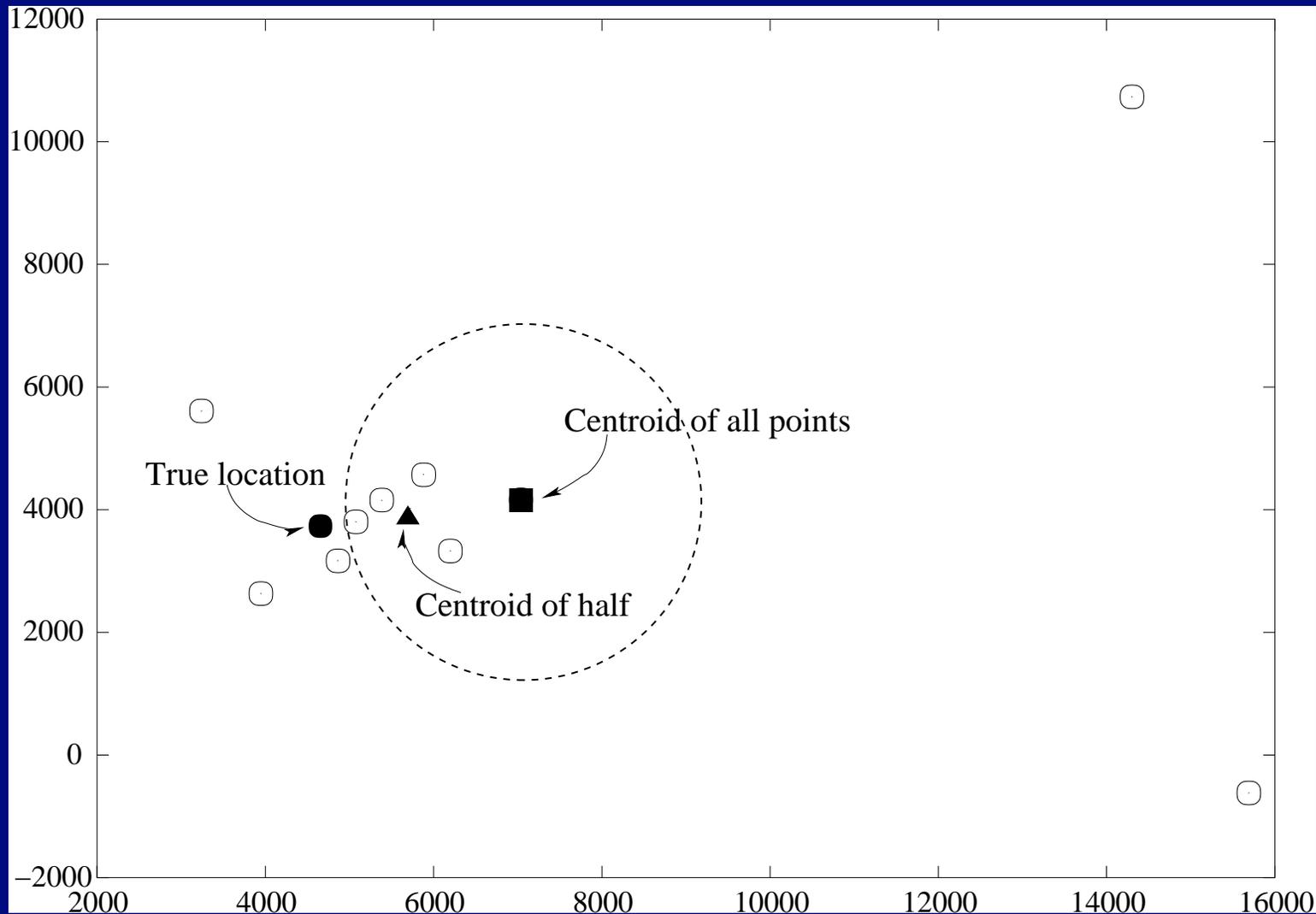
error control - propagation



error control - small angles



error control - remove outliers

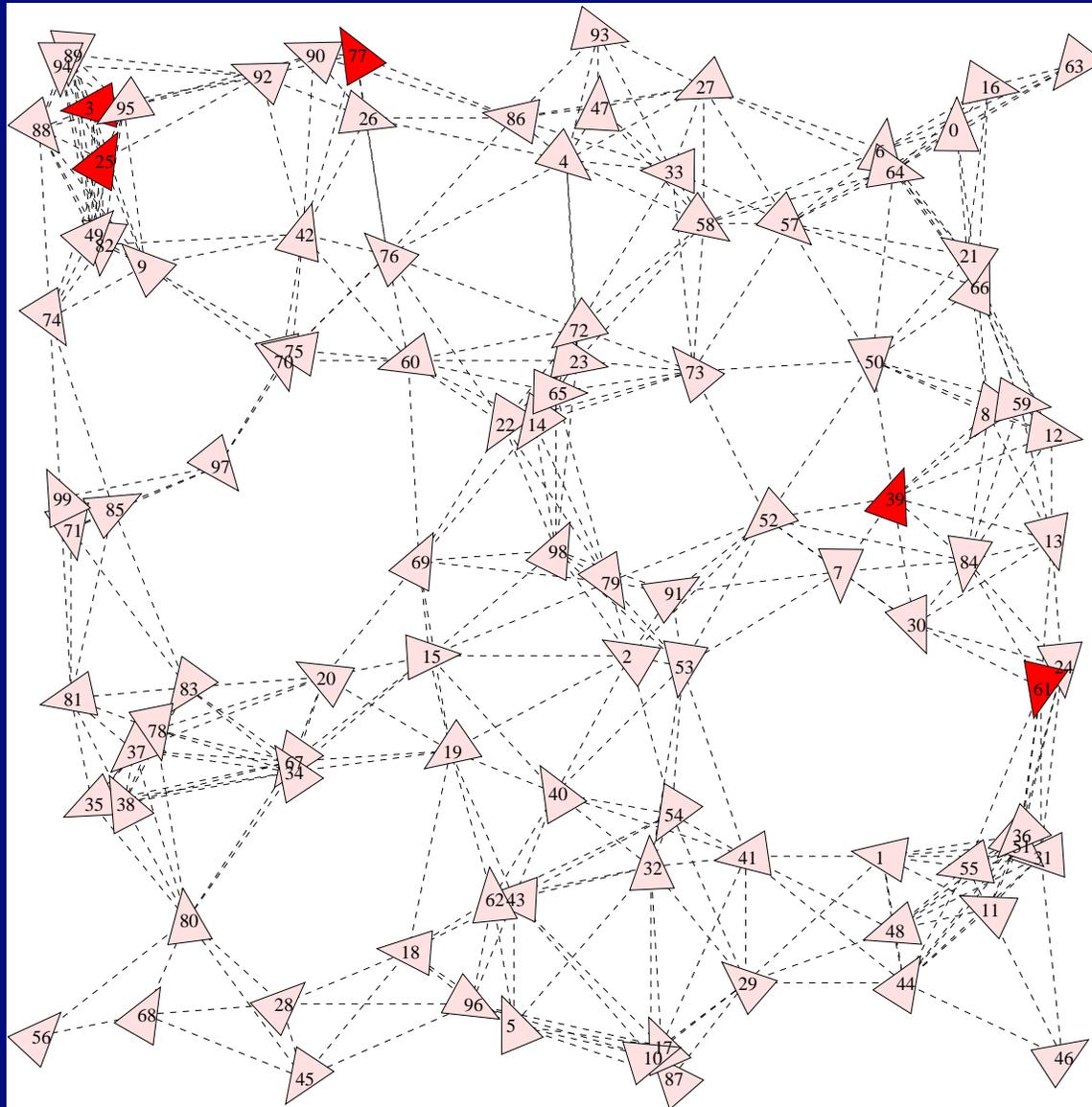


simulation

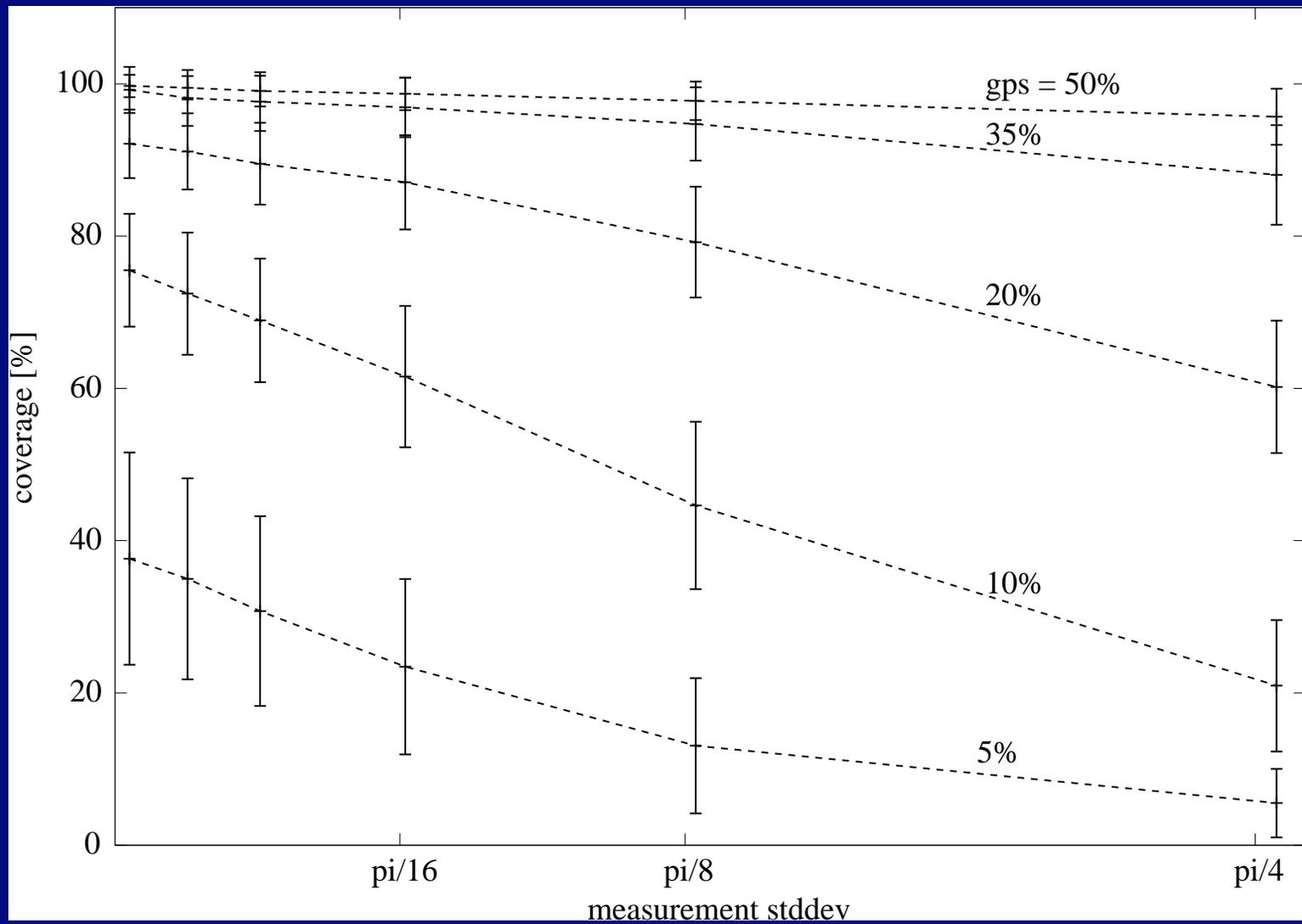
- random topology 100 nodes
 - isotropic ¹
 - AoA errors
 - white Gaussian noise
 - with a 95% probability, the measured angle is within $\pm 2 \textit{stddev}$ of mean

- performance metrics
 - coverage - how many nodes get a position/heading
 - absolute position error - in number of hops
 - bearing error - with respect to landmarks
 - heading error - absolute orientation

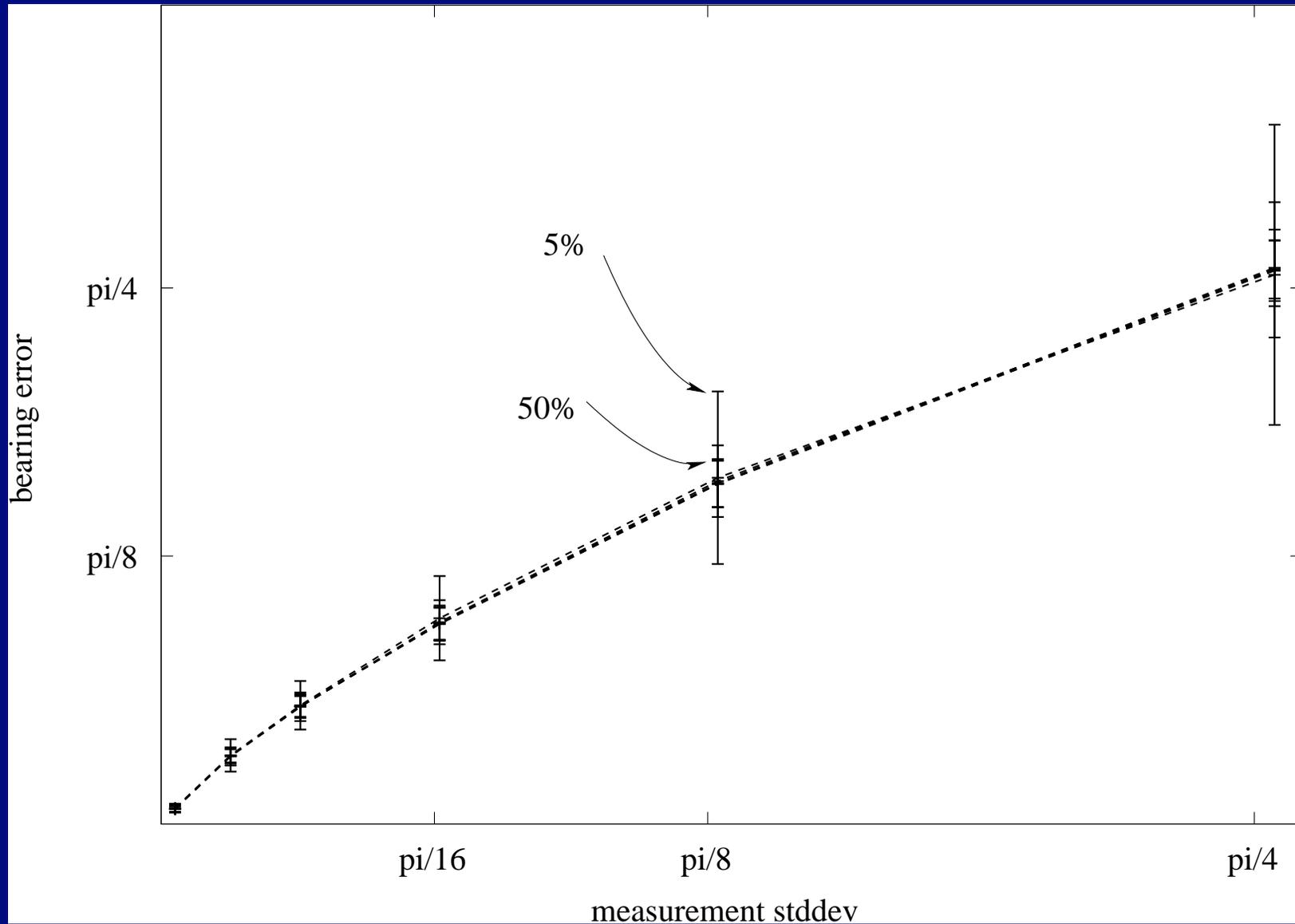
¹the network has the same properties (density, radio range) in all directions



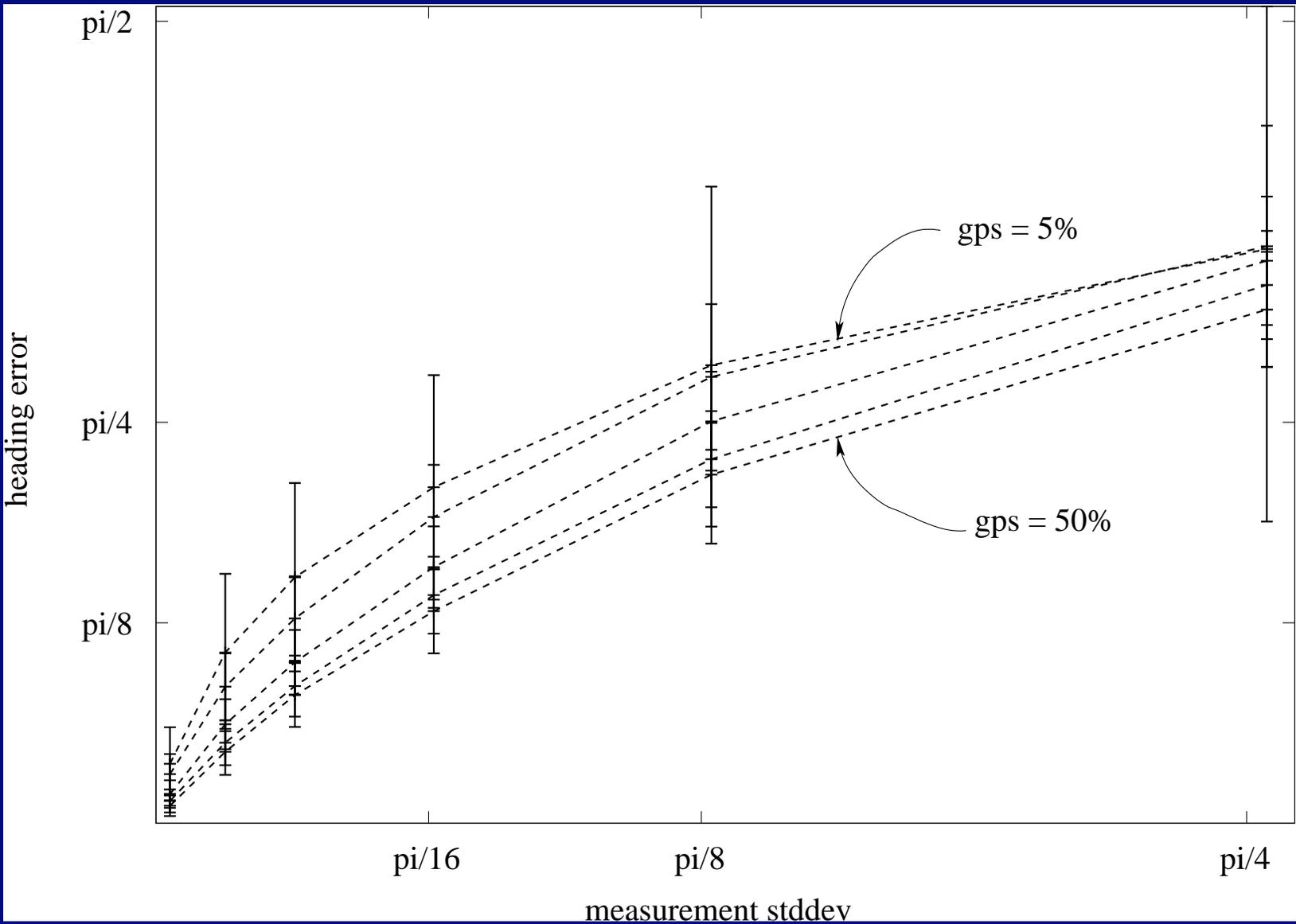
coverage



bearing error

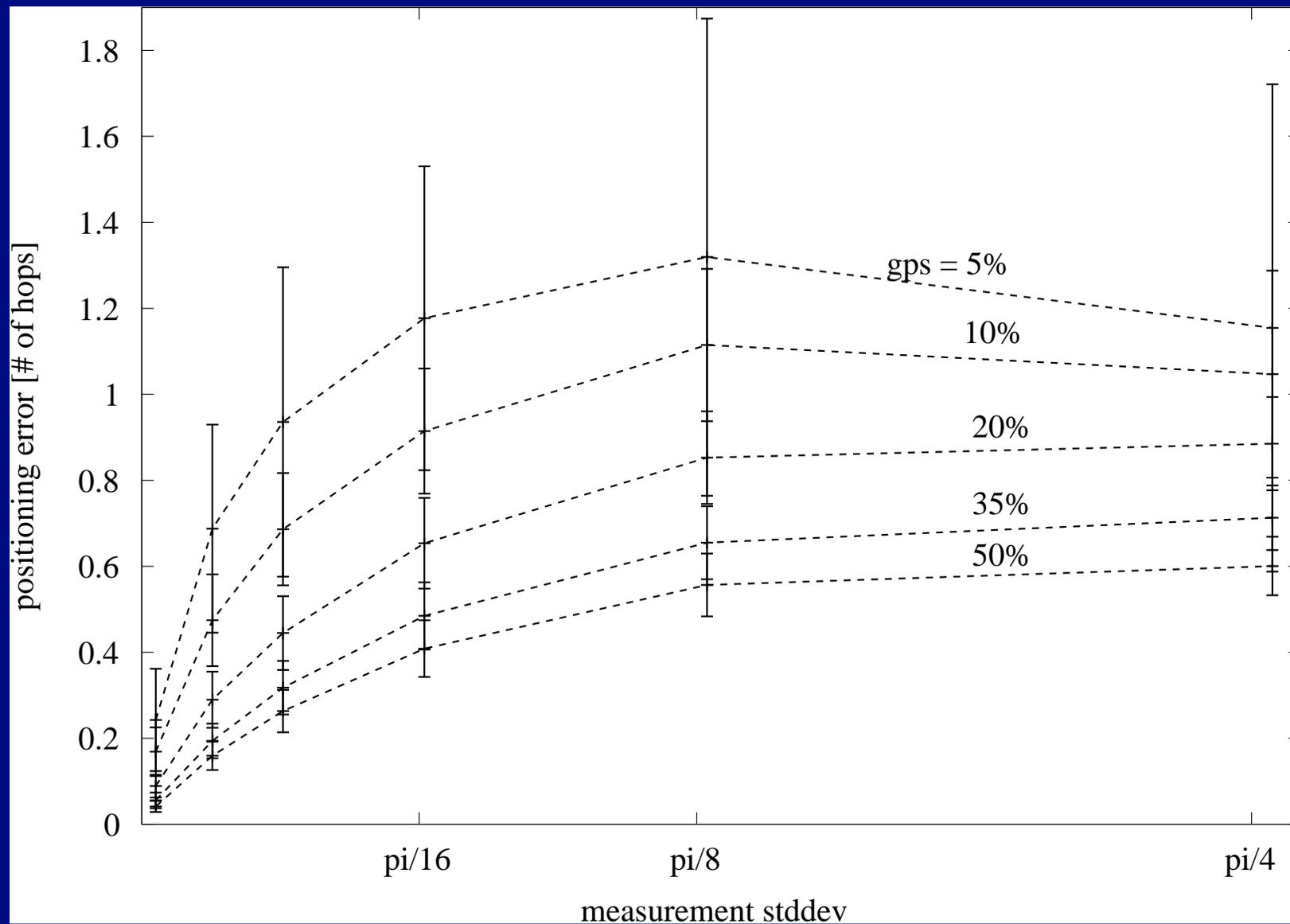


heading error



○ Heading error = double of the bearing error

location error



simulation summary

- error reduction - simple methods
- error - coverage tradeoff
- heading error is about double the heading error
- nodes inside the convex hull of acquired landmarks get better estimates
- # of landmarks is more critical for position than for orientation

future work

○ multimodal estimation

- can AoA and signal strength be used together?

○ node mobility

- a moving landmark
 - is a new landmark
 - one flying landmark could be enough for the entire static network
- mobile nodes are supported by static nodes

conclusions

- APOS (Ad Hoc Positioning and Orientation System)
 - provides position and orientation for randomly deployed nodes
 - needs AoA capability in all nodes, but no signal strength
 - distributed, no infrastructure
 - uses a DV based scheme to propagate landmark bearings
 - positioning accuracy → one hop away from the true location
 - orientation accuracy → double the measuring accuracy